

Exploration with Limited Memory: Streaming Algorithms for Coin Tossing, Noisy Comparisons, and Multi-Armed Bandits

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Abstract

- Finding the most biased coin by tossing -- a classical exploration problem in computer science and machine learning.
- Assuming a gap parameter Δ , elimination-based algorithms have provided solution with $O(\frac{n}{\Delta^2})$ coin tosses which matches the lower bound.
- However, these algorithms inherently require storing all the coins, which is not memory-efficient.
- We studied the sample-space trade-off under the streaming coin tossing mode: the algorithm can only toss an incoming or stored coin.
- We designed an algorithm which only stores a **single extra coin**, which means the sample-space trade-off does not exist.
- En route to the one-coin algorithm, we also proposed preliminary memory-efficient algorithms with $O(\log(n))$, $O(\log\log(n))$ and $O(\log^*(n))$ stored coins.
- Extensions of our main algorithm includes finding the k most biased coins and other exploration problems. E.g.. Finding top- k elements using noisy comparisons; Finding an ϵ -best arm in stochastic multi-armed bandits.

Background

Goal: Find the **most biased coin (denote as coin*)** with **high constant probability** by tossing each coin several times.

Parameters:

- n The number of coins;
- Δ The gap between the most and second-most biased coins

A good algorithm should:

- ✓ Return the most biased coin w.h.p. (**correctness**)
- ✓ Use small number of tosses (**sample complexity**)
- ✓ Store small number of coins (**space complexity**)

A Naïve Algorithm:

- Toss each coin $O(\frac{\log(n)}{\Delta^2})$ times
- Sample complexity $O(\frac{n}{\Delta^2} \log(n))$
- Space complexity: 1 coin

Median Elimination Algorithm (Even-Dar et al., 2006):

- Round 1: toss each coin $O(\frac{1}{\Delta^2})$ times
- Round 2+: eliminate $\frac{1}{2}$ of the coins; increase 1.5x coin tosses
- Sample complexity $O(\frac{n}{\Delta^2})$
- Space complexity: $\Omega(n)$ coins

Any sample-space trade-off?

-- streaming model: To toss a past coin, must have stored it.

Our Contribution

Main Theorem (Assadi and Wang, 2020)

There exists a streaming algorithm that given n coins arriving in a stream with the gap parameter Δ and confidence parameter δ , finds the most biased coin with probability at least $1 - \delta$ using $O(\frac{n}{\Delta^2} \cdot \log(1/\delta))$ coin tosses and a memory of a single coin.

- No sample-space trade-off!
- Preliminary : $O(\log(n))$, $O(\log\log(n))$ and $O(\log^*(n))$ coins memory algorithms.
- Additional result: Top- k coin exploration with $O(k)$ coin memory.
- Additional results: Noisy comparisons and Multi-Armed Bandits.

Preliminary Algorithms

The $O(\log(n))$ -Coin memory Algorithm:

- **Multiple levels:** 4-coin memory per level
- **Level 1:** toss each coin $\frac{30}{\Delta^2}$ times; send the **most biased** to the level 2.
- **Level 2+:** increase the number of tosses by 1.5x
- ✓ **Correctness:** Probability of losing coin* exponentially decreases.
- ✓ **Sample complexity:** i -th level: $\frac{30}{\Delta^2} \cdot (1.5)^{i-1} \cdot \frac{n}{2^{i-1}}$, overall $O(\frac{n}{\Delta^2})$.
- ✓ **Space Complexity:** $O(\log(n))$ levels; each level 4 coins.

The $O(\log\log(n))$ and $O(\log^*(n))$ Coin Algorithms:

- $O(\log\log(n))$ memory: stopping at the $\log\log(n)$ level
- $O(\log^*(n))$ memory: aggressive selections of coins (iterative logarithm factor) and increments of coin tosses (tower factor) (cf. [Agarwal et al., 2017])

Main Algorithm – One Coin Suffices

Idea:

- Pick **only one coin** to store, name as *King*.
- Worst case $\Theta(n)$ coins challenge the *King* -- give the *King* privilege: **only be dethroned if lost multiple levels of challenge**.
- Bound the **sample complexity**: limit the tosses of the *King* by **budget**.

Algorithm GAME-OF-COINS:

- For **each arriving coin** give the *King* a **budget** of $O(\frac{1}{\Delta^2})$.
- To challenge the *King*, toss both coins $\frac{30}{\Delta^2} \cdot (3)^{i-1}$ times at level i ;
- A *King* is defeated **only if it exhausts all its budget**.

Analysis:

- ✓ **Sample Complexity:** At most $2n \cdot O(\frac{1}{\Delta^2})$ budgets $\rightarrow O(\frac{n}{\Delta^2})$ coin tosses.
- ✓ **Space Complexity:** Only store 1 coin.
- ✓ **Correctness:**
 1. The coin* can exhaust the budget of other *King* (**soundness**)
 2. If coin* as the *King* \rightarrow budget sufficient *in expectation*.
 3. Control the variance:
 - a) The budget behaves like random walks (but with flexible length).
 - b) The challenging rule \rightarrow budget distribution sub-exponential.
 - c) Beating the union bound by Bernstein inequality (**completeness**).

Extensions of the Algorithm

Algorithm for top- k coins:

- Main technical contribution -- a delayed challenging rule & a potential function argument.
- Avoid eliminating any top- k coin -- use a buffer to swap defeated coins (**correctness**).
- Number of coins eventually decreases -- bounded **sample complexity**.

Noisy Comparisons and ϵ -PAC Multi-Armed Bandit (MAB):

- Noisy comparison – $O(k)$ space algorithm for finding top- k elements.
- No gap guarantee -- a $O(\log^*(n))$ space algorithm. Most recently, an extension to a 2-arm algorithm.

Extensions and open problems:

- The **instance-sensitive sample complexity**: $H_2 = O(\sum_{i>1} \frac{1}{\Delta_i^2} \log \log(\frac{1}{\Delta_i}))$.
- Single-pass: achievable with random arrival of coins and a value $O(H_2)$.
- Single-pass with lower bounds; arbitrary stream with $O(\log(\frac{1}{\Delta^2}))$ passes [Jin et al., 2021].
- **Open:** tight number of passes to achieve $O(H_2)$ sample complexity.



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