



# Exponentially Improved Dimensionality Reduction for $\ell_1$ : Subspace Embeddings and Independence Testing

Yi Li<sup>1</sup>, David P. Woodruff<sup>2</sup>, Taisuke Yasuda<sup>2</sup>  
Nanyang Technological University<sup>1</sup>, Carnegie Mellon University<sup>2</sup>

## Overview

A celebrated result of Johnson and Lindenstrauss [JL84] gives highly efficient dimension reduction for the  $\ell_2$  norm. That is, a  $r \times n$  Gaussian matrix  $S$  for  $r = \Theta(\varepsilon^{-2} \log \delta^{-1})$  satisfies  $\|Sx\|_2 = (1 \pm \varepsilon)\|x\|_2$  with probability at least  $1 - \delta$ . However, the  $\ell_1$  norm is often more appropriate than the  $\ell_2$ :

- For data mining applications, the  $\ell_1$  norm is more robust to outliers than the  $\ell_2$  norm
  - The  $\ell_1$  norm is twice the total variance distance for distributions
- We thus seek analogous results to the Johnson-Lindenstrauss lemma for the  $\ell_1$  norm.

## Subspace Embeddings

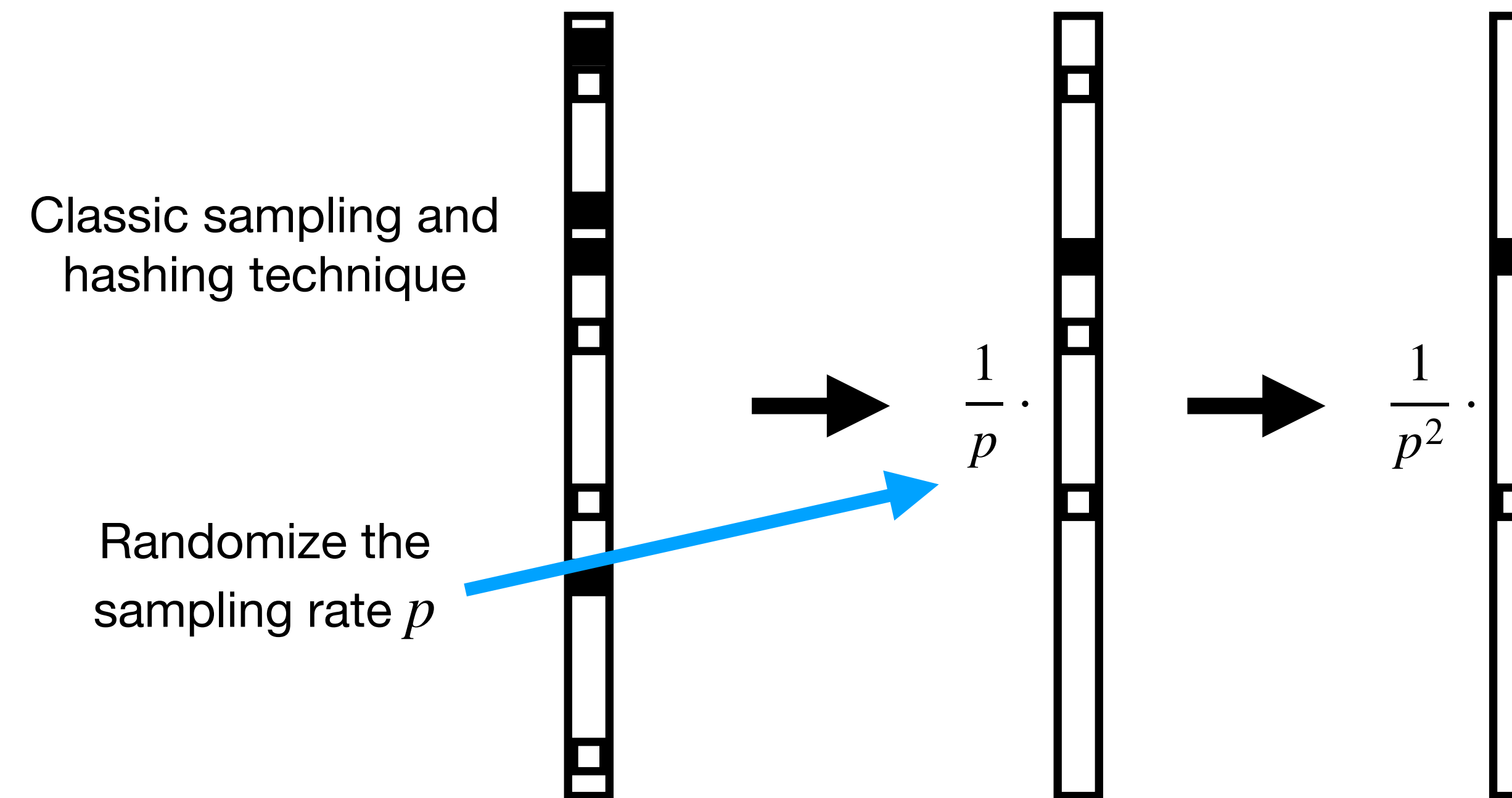
The known  $\ell_1$  analogue for the Johnson–Lindenstrauss lemma, i.e. linear oblivious dimension reduction maps  $S$  with  $\|Sx\|_1 = (1 \pm \varepsilon)\|x\|_1$ , has a doubly exponential dependency on  $\varepsilon^{-1}$  [WW19]. For subspace embeddings, i.e. the above guarantee for every vector in a  $d$ -dimensional subspace, [WW19] also achieves a doubly exponential dependence on  $d, \varepsilon^{-1}$ , and also shows a singly exponential lower bound in  $d$ . Can we close this gap? **We answer this question in the affirmative.**

	$\ell_2$ [JL84]	$\ell_1$ LB [WW19]	$\ell_1$ UB [WW19]	$\ell_1$ UB [LWY21]
1 vector	$\varepsilon^{-2}$		$2^{2^{\varepsilon^{-2}}}$	$2^{\varepsilon^{-1}}$
$m$ vectors	$\varepsilon^{-2} \log m$	$2^{\sqrt{m}}$	$2^{2^{\varepsilon^{-2} \log m}}$	$2^{\varepsilon^{-1} m}$
$d$ -dim subspace	$\varepsilon^{-2} d$	$2^{\sqrt{d}}$	$2^{2^{\varepsilon^{-2} d}}$	$2^{\varepsilon^{-1} d}$

\*Suppresses big Oh and log factors

## Idea for Singly Exponential Dependence in $\varepsilon^{-1}$

To achieve our singly exponential dependence on  $\varepsilon^{-1}$ , we start with the M-sketchn construction of [CW14]. M-sketchn is based on sampling and hashing the coordinates of  $x$ , and achieves a distortion of  $O(1)$ . We modify this construction by **randomizing the sampling rates** themselves to achieve our result.



## Idea for Singly Exponential Dependence in $d$

To achieve our singly exponential dependence on  $d$ , we cannot afford to union bound over a net, as done by [JL84]. Instead, we apply our earlier result on the  $\ell_1$  **leverage score vector** with distortion  $(1 + \varepsilon/d)$ .

## Clarkson–Drineas–Magdon-Ismail–Mahoney–Meng–Woodruff (2013)

Let  $A \in \mathbb{R}^{n \times d}$ . Then, there exists a vector  $\lambda = \lambda(A) \in \mathbb{R}^n$  such that  $\|\lambda\|_1 = 1$  and for all  $x \in \text{span}(A)$ ,

$$\frac{|x_i|}{\|x\|_1} \leq d \cdot \lambda_i$$

for every  $i \in [n]$ .

## Independence Testing

Consider a distribution given by a data stream:

- Each stream element is  $(i_1, \dots, i_q)$  for  $i_j \in [d]$

- **Empirical joint distribution  $P$ :**

$$p(i_1, \dots, i_q) = \frac{\text{number of occurrences of } (i_1, \dots, i_q)}{\text{length of stream}}$$

- **Empirical product distribution  $Q = Q_1 \times \dots \times Q_q$ :**

$$q_j(i) = \frac{\text{number of occurrences of } (*, \dots, *, i, *, \dots, *)}{\text{length of stream}}$$

Our task is to estimate  $\|P - Q\|_1$ .

The previous known algorithm for this problem has a doubly exponential dependence on  $q$ :

## Braverman–Ostrovsky (2010)

$\|P - Q\|_1$  can be estimated in  $(\varepsilon^{-1} \log d)^{q^{O(q)}}$  space.

We improve this to a singly exponential bound:

## Li–Woodruff–Yasuda (2021)

$\|P - Q\|_1$  can be estimated in  $2^{O(q^2)} (q\varepsilon^{-1} \log d)^{O(q)}$  space.

## Our Construction

We design a sketching matrix  $S$  which takes the tensor product structure  $S = (S^1) \otimes \dots \otimes (S^q)$ , where each  $S^j$  sketches each mode of the tensor. This allows us to maintain  $S^1 Q_1, \dots, S^q Q_q$  in the stream and compute  $SQ = (S^1 Q_1) \otimes \dots \otimes (S^q Q_q)$ . We also maintain  $SP$ , then estimate  $\|P - Q\|_1$  based on  $SP - SQ = S(P - Q)$ .

As with the subspace embedding, our sketch construction for a single mode starts with sampling and hashing techniques.

This is used recursively to handle all  $q$  modes.

## References

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